# Interconnection and Competition among Asymmetric Networks in the Internet Backbone Market

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### Abstract

We examine the interrelation between interconnection and competition in the internet backbone market. Networks asymmetric in size choose among different interconnection regimes and compete for end-users. We show that a direct interconnection regime, Peering, softens competition compared to indirect interconnection since asymmetries become less influential when networks peer. If interconnection fees are paid, the smaller network pays the larger one. Sufficiently symmetric networks enter a Peering agreement while others use an intermediary network for exchanging traffic. This is in line with considerations of a non-US policy maker. In contrast, US policy makers prefer Peerings among relatively asymmetric networks.

Keywords: Internet Backbone, Endogenous Network Interconnection, Asymmetric Networks, Two-Way Access Pricing

JEL Classification: L10, L96, D43

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### 1 Introduction

The internet as a communications industry is subject to network externalities. These externalities have forced Internet Backbone Providers (IBPs) to interconnect with each other in order to provide their customers with "world-wide connectivity", hence increasing consumers' benefits and willingness-to-pay for internet access. From an economic perspective there are several ways to interconnect with other networks. The specific type of interconnection influences competition for end-users, and vice versa.

This paper aims to provide a general analysis of the industrial organization of an unregulated internet backbone market. We endogenize both networks' interconnection and competition decisions while explicitly accounting for asymmetric network sizes, which are widely observed in practice. We will study the following questions: What determines networks' choice of interconnection? How do different types of interconnection affect competition for end-users? Who pays whom for interconnecting networks? Are networks' decisions in line with welfare considerations?

We will suggest to consider a new interconnection regime, Paid Peering, and find that networks which are sufficiently symmetric in size prefer it (together with better known Bill-and-Keep Peering) over using an intermediary network to exchange data. For medium ranges of network asymmetry, Paid Peering even dominates both alternative interconnection regimes: networks can raise profits in comparison to a situation where they were restricted to the choice of Bill-and-Keep Peering versus IP-Transit. Only for large asymmetries, they buy IP-Transit from an intermediary network in equilibrium. Our model will suggest that this interconnection behavior is not always desirable from a welfare point of view. Finally, taking into account that the market for IP-Transit is dominated by US carriers, a non-US trade policy oriented regulator would find that there is too much Peering and would seek to restrict Peering of networks which are sufficiently asymmetric in size.

Our model has the following timing: first, two networks, which are ex ante connected via an intermediary backbone, negotiate their interconnection regime. In case of Paid Peering, they bargain for a settlement-fee (interconnection fee or access price) that could flow either direction on stage two of the game. Third, they compete in prices for consumers with heterogeneous preferences in a Hotelling model. Finally, consumers choose the network maximizing their net benefits.

Our results are strongly driven by asymmetries in network sizes: the larger the ex ante asymmetry is, the larger the profit differences between the networks when using an intermediary backbone, which in turn serve as threat points in the Nash bargaining game. As a consequence, both the settlement-fee resulting from the bargaining process and the interconnection decision reached in equilibrium depend on the degree of network asymmetry.

There is a rapidly developing literature on interconnection and two-way access pricing in telecommunications, which one might think of being related to the internet backbone market. Armstrong (1998) and Laffont et al. (1998) constitute two fundamental works, while Vogelsang (2003) provides a comprehensive survey on this literature. However, there are two crucial differences which make an adoption of the analysis on the telecommunications market to the internet backbone highly problematic: First, interconnection in the internet backbone is not subject to regulation. Cash flows associated with interconnection on the internet do not depend on the direction of traffic but may be negotiated freely in the market.<sup>1</sup> Second, destination based price discrimination is usual in telecommunications, while it is practically impossible on the internet.<sup>2</sup>

There is also a more recent theoretical literature on telecommunications relaxing these industry specific restrictions: Carter and Wright (2003), Armstrong (2004), Gilo and Spiegel (2004) and Peitz (2005) study competitively chosen asymmetric access prices, asymmetric networks or IP-Transit as an outside option when negotiating the terms of interconnection. Yet our paper is the first to unify all three issues in one model.

Focusing on the internet, Laffont et al. (2003) study the strategic behavior of backbone operators in an environment of reciprocal access pricing in two-sided markets. Mendelson and Shneorson (2003) extend this framework to consumer delay costs and capacity decisions. Contrarily, because of already existing world-wide

<sup>&</sup>lt;sup>1</sup>In telecommunications, regularly the data sending network has to pay the receiving network for terminating a call. Moreover, policy makers often require such termination charges or "access charges" to be set reciprocally.

<sup>&</sup>lt;sup>2</sup>It is standard for consumers to pay more for long-distance or international phone calls than for local ones. To imitate such price discrimination on the internet, a consumer would have to be asked before each click on a Web link whether she would be willing to pay a specific price depending on the network distance to a specific target Web site's location.

connectivity we abstract from network externalities in consumers' utility functions. Because of the unregulated nature of the internet backbone, we let networks negotiate access prices freely.

Crémer et al. (2000) analyze in a Cournot model (thus endogenizing capacity) whether dominant network operators have incentives to lower the interconnection quality to rival networks. By extending the Katz and Shapiro (1985) network competition model they show that a network with a large installed base of customers is likely to degrade its interconnection quality with smaller networks.<sup>3</sup> However, nowadays there is excess capacity all over the backbone market,<sup>4</sup> and the marginal costs of data transmission are virtually zero.<sup>5</sup> Therefore, instead of modelling competition based on capacities/quantities we focus on price competition with differentiated products in the retail market against the background of (exogenous) competition in the Transit market. Instead of competition based on quality of interconnection we assume perfect transmission quality, which is due to existing world-wide connectivity and the absence or bottlenecks, and let networks choose among several interconnection regimes.<sup>6</sup>

The papers connected most closely to our's are Baake and Wichmann (1999) and Besen et al. (2001) in the sense that they also endogenize the choice of IBPs' interconnection regime. The former studies the Transit vs. Peering decision in the context of quality differentials though, while the latter provides a bargaining process of Peering partners (implicitly introducing the option for Paid Peering). Both do not consider effects on competition for end-users. To the best of our knowledge, our paper is the first attempt to endogenize both networks' interconnection and compe-

<sup>&</sup>lt;sup>3</sup>Foros and Hansen (2001) also study interconnection quality and competition between IBPs but derive opposing results concerning the development of market shares. Roson (2002) provides a more thorough discussion of Crémer et al. (2000) and that article. Foros et al. (2005) analyze interconnection in a two-stage game where networks first decide about interconnection quality and compete in quantities thereafter.

<sup>&</sup>lt;sup>4</sup>Telegeography, a consultancy, notes: "Despite significant and consistent growth in data traffic flows across the world's communications networks, a huge portion of potential network capacity remains unused. [...] only three percent of the maximum possible intercity bandwidth in Europe and the U.S. has been 'lit' for service provision." (http://www.telegeography.com/press/releases/2005-04-20.php)

<sup>&</sup>lt;sup>5</sup>See Atkinson and Barnekov (2004) or Nuechterlein and Weiser (2005, p.38).

<sup>&</sup>lt;sup>6</sup>In the internet backbone, excess capacity leads to virtually perfect quality of interconnection.

tition decisions among asymmetric networks while taking into account the economic differences between the internet backbone and telecommunications markets.

The paper is organized as follows. Section 2 describes the most widely used interconnection regimes in more detail. Section 3 sets the stage for the model and derives networks' equilibrium prices, market shares and profits under Intermediary and Bill-and-Keep Peering regimes, respectively. Section 4 introduces the Nash bargaining game used under the Paid Peering regime. Section 5 examines incentives to peer and defines parameter ranges of interconnection equilibria. Section 6 takes a welfare perspective. Section 7 offers an outlook on the internet backbone market's future, while section 8 concludes.

# 2 Interconnection Practice in the Internet Backbone Market

The market for interconnection between network operators on the internet has developed rapidly during the last years and is expected to do so in the future, too. According to one forecast (IDC, 2003), the volume of global internet traffic should nearly double annually, increasing from 180 petabits per day in 2002 to 5,175 petabits per day by the end of 2007.

Up to the mid-1990's, large parts of the internet were owned by governmental agencies, for instance the National Science Foundation (NSF) in the US. Along with deregulating telecommunications markets those agencies steadily withdrew from managing the infrastructure, which was followed by a rapid technological and commercial development. As is common knowledge, the internet itself is a network of networks playing an intermediary role in connecting consumers who are connected to it via Internet Service Providers (ISPs). Traditionally, ISPs have formed regional networks that exchange traffic generated by their customers via Internet Backbone Providers. Nowadays, however, more and more ISPs are vertically integrated into IBPs.<sup>7</sup>

Users could be assigned to two segments, traffic senders (e.g. Web sites) and traffic receivers (e.g. end-users). But due to the emergence of new broadband

<sup>&</sup>lt;sup>7</sup>Kende (2000) and Atkinson and Barnekov (2004) provide non-formal studies of the internet backbones' market environment pointing on currently important issues and open questions.

services, for instance Voice over IP (VoIP) or video conferencing, the distinction blurs and consumers become senders and receivers at the same time. Therefore, the direction of traffic flows cannot be attributed clearly to distinct utility levels of senders and receivers, respectively.<sup>8</sup>

How does traffic get from consumer 1 to consumer 2? Suppose 1 and 2 are video conferencing and are connected to different networks. Those have two main options to exchange traffic, Transit and Peering.

IP-Transit/Intermediary: If a direct connection is not feasible or desirable, two networks can buy so-called Transit services from a third network. Under such an arrangement both networks pay a variable charge per unit of traffic to the intermediary network which obligates to deliver the traffic to any specified destination and from a certain origination. For being able to fulfil this obligation, networks offering Transit mostly have a large physical network and are connected to many other networks via Peering or further Transit sales. The IP-Transit market is dominated by so-called Tier-1 networks which are mainly US based.<sup>9</sup>

Peering: <sup>10</sup> Bill-and-Keep Peering, also called settlement-free Peering, has evolved as the regular type of direct interconnection regime between two networks since privatization. Networks exchange traffic without charging any fees to each other. However, under such a Peering agreement no participating network has the obligation to terminate traffic to or from a third party. Each network must only process traffic from the Peering partner to its own customers (and the customers of their customers and so on), but not to the remainder of the internet. This constitutes a major difference between IP-Transit and Peering. In our example, customers 1 and 2 can exchange traffic without causing any interconnection costs to the networks they have subscribed to if those are peering. A necessary requirement for Peering in general is the exact routing of traffic in order to control the flow of traffic. Oth-

<sup>&</sup>lt;sup>8</sup>A widely-used billing mechanism for Transit services is the so-called 95th Percentile Billing which does only account for traffic volume but not for traffic direction.

<sup>&</sup>lt;sup>9</sup>A network is regarded to have Tier-1 status if it is connected to the whole internet while never paying for interconnection itself.

<sup>&</sup>lt;sup>10</sup>In the industry, there is a difference between "Private Peering", where exactly two networks build or lease lines to interconnect, and "Public Peering" where several networks interconnect their lines in a node, a so-called Internet Exchange Point. As economic differences are not very significant and more and more networks use Private Peering, we only consider this type in our model. See Kende (2000) for more details.

erwise, a third network which has been denied Peering by one of the participating networks could free-ride on the existing Peering arrangement.

A Paid Peering regime between two networks implies the same rights concerning their exchange of traffic. In contrast to a Bill-and-Keep arrangement, one network charges the other for exchanging traffic. We may emphasize that Paid Peering is a relatively new type of interconnection regime and has only recently begun to be employed. In our example, suppose the network of customer 1 agreed to pay for traffic exchange with the network 2 belongs to, thereby forming a Paid Peering interconnection regime. In this environment it has no impact on the stream of money whether 1 sends an e-mail to 2, or vice versa: in both cases 1's network will pay 2's. However, since Paid Peering is no Transit contract, 2's network will not proceed traffic from 1 to a third party which is not a customer of 2's network.

# 3 The Model

### 3.1 Key Assumptions

There are two networks  $i \in \{A, B\}$  each having a fixed installed base of customers  $\alpha_i$  that is not subject to competition.<sup>13</sup> Without loss of generality we assume  $\alpha_A \geq \alpha_B$ . On top,  $\bar{\alpha}$  consumers are situated in a *battlezone*, where networks A and B compete in prices.<sup>14</sup> We assume excess capacity on the part of the networks so they could serve battlezone consumers without extra investments. This reflects the current infrastructure environment. Ex ante both networks are connected to the remainder

<sup>&</sup>lt;sup>11</sup>A Paid Peering settlement could appear in several different forms of payment, either fixed amount payments or a variable charge per unit of traffic (or a combination of both). The previous type of payment could involve asymmetric cost sharing regarding the technological fixed costs of installing traffic exchange points between Peering partners.

<sup>&</sup>lt;sup>12</sup>For more details on internet traffic, see Giovanetti and Ristuccia (2005) or Kende (2000).

<sup>&</sup>lt;sup>13</sup>Internet Service Providers selling internet access to those consumers are integrated into the networks.

 $<sup>^{14}</sup>$ The most intuitive explanation is of a geographic nature: the locked customers of network i can only be connected directly to network j for prohibitively high costs, e.g. because they live in a rural area. The battlezone, however, consists of consumers living in large cities where both networks have a point of presence (POP). Another interpretation is that A and B compete in new services, e.g. Voice over IP, in the battlezone but also have legacy customers who are not interested in such services.

of the internet by using an intermediary, thereby offering their customers world-wide connectivity. As there is Bertrand price competition in the market for IP-Transit, we assume the intermediary to be the cheapest tier-one network available, by definition offering access to all remaining consumers connected to the internet,  $\kappa$ . It is not important whether the intermediary directly serves the  $\kappa$  consumers as ISP or connects other networks' consumers via its backbone to networks A and B. There is a continuum of consumers, of mass 1, so  $\alpha_A + \alpha_B + \bar{\alpha} + \kappa = 1$ . Figure 1 charts the competition set-up.

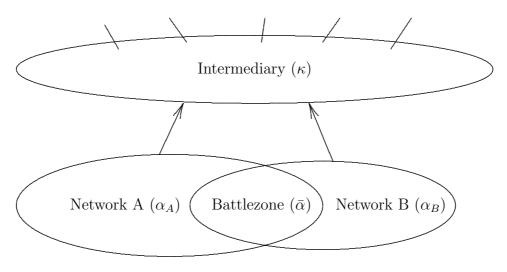


Figure 1: Network interconnection via an intermediary

Networks' cost structure:

• Networks face an exogenous market price for upstream Transit,  $t_u$ , per unit of data

<sup>&</sup>lt;sup>15</sup>In line with this, we model no quality differentials among Peering and Transit, unlike as in Crémer et al. (2000) or Baake and Wichmann (1998), since, according to industry representatives, there is no clear relationship between interconnection quality and regimes. Consequently, demand-side network effects do not play a role in the model since customers enjoy world-wide connectivity on a constant quality level regardless of the networks' interconnection decision or competition.

<sup>&</sup>lt;sup>16</sup>In our model we do not cover competition where one of the two networks has tier-one status. Therefore, we do not endogenize the intermediary's price of IP-Transit. See Prüfer and Jahn (2005) for a discussion of the influence of Bertrand competition on the internet backbone industry's outlook and market structure.

- Technical marginal cost of sending data are zero. We discuss the theoretical implications of this fact on the price for IP-Transit in section 7.
- Costs of connecting customers to a network within the battlezone are symmetric and, for simplicity, normalized to zero.
- In case of a Peering arrangement, each network has to bear a fixed cost F, where  $0 < F \le t_u \bar{\alpha} (\alpha_A + \alpha_B + \frac{\bar{\alpha}}{2}) + 2\alpha_A \alpha_B t_u \equiv t_u w.^{17}$

Since top-level backbones do not charge different fees for upstream or downstream traffic, we merely assume that each consumer sends one unit of data to each other consumer and receives one unit of data from each other consumer, thereby not taking into account which network the other consumer is connected to (balanced calling pattern). This yields each consumer a gross benefit, v. Finally, we assume that prices  $p_i^L$  in the locked areas are not affected by competition in the battlezone where both networks charge each customer a price  $p_i$ . In case, network i is a monopolist in its locked area, assume the level of  $p_i^L$  equals a price cap set by a regulatory body.<sup>18</sup>

### 3.2 The Game

The timing of the game is as follows:

- 1. Networks A and B decide non-cooperatively about the interconnection regime between them, *Intermediary*, *Bill-and-Keep Peering* or *Paid Peering*. If networks cannot agree on a specific Peering regime, both are forced to use the intermediary.
- 2. In case of Paid Peering, networks bargain for a fixed settlement which could flow either direction.
- 3. Networks A and B set prices  $p_i$  for consumers in the battlezone. Consumers have heterogeneous preferences, so networks compete in a Hotelling-like environment. Heterogeneity could depend on different complementary services

 $<sup>^{17}</sup>F$  encompasses all fixed-step costs for setting up a physical interconnection, buying routers, etc. and organizational costs for managing a Peering agreement. Without an upper boundary for F,  $t_u w$ , Peering can never be an equilibrium, which makes the analysis less interesting.

<sup>&</sup>lt;sup>18</sup>Hence we have  $p_i \leq p_i^L$  where marginal cost are zero.

offered by the networks, e.g. specific Web content or software applications certain consumers are already used to. Note that heterogeneity refers to the retail market of internet access, while data exchange between networks is a homogenous good.

4. Consumers in the battlezone choose the network maximizing their net benefits.

We will derive equilibrium profits of both networks under Bill-and Keep Peering (BK) and Intermediary regimes at the third stage of the game first, derive Paid Peering (PP) profits at the second stage and compare them at the first stage afterwards to yield incentives for choosing the regimes. Then we derive parameter constellations where choosing Bill-and-Keep Peering, Paid Peering or Intermediary constitute equilibrium strategies for both networks.

### 3.3 Price Competition under the Intermediary Regime

Consider a standard Hotelling (1929) model. Consumers are indexed by x and uniformly distributed on the interval [0,1] with increasing preference for network B. The network differentiation parameter (transportation cost parameter) is  $\tau > 0$ , so that a consumer's utility function is given by

$$U = \begin{cases} v - \tau x - p_A & \text{if buying from network A} \\ v - \tau (1 - x) - p_B & \text{if buying from network B} \\ 0 & \text{otherwise.} \end{cases}$$
 (1)

We assume  $v \geq \frac{3}{2}\tau + 2\kappa t_u$  to assure the market is covered. It is simple to calculate the standard marginal consumer who is indifferent between A and B and denoted by

$$\hat{x} = \frac{1}{2} + \frac{p_B - p_A}{2\tau}. (2)$$

Note that  $\hat{x}$  also specifies A's market share within the battlezone, while  $(1 - \hat{x})$  is B's battlezone market share. Profit functions<sup>19</sup> under the Intermediary regime are

<sup>&</sup>lt;sup>19</sup>We assume that networks are able to discriminate prices between locked consumers and the battlezone. If that was not possible, as  $\alpha_A \ge \alpha_B$  there would be no price Nash equilibrium in pure strategies. Therefore, and because we believe in the feasibility of price discrimination based on the sender's—not the receiver's—location in the internet, we restrict our analysis to this case.

given by

$$\Pi_A^I = \hat{x}\bar{\alpha}(p_A - 2\kappa t_u) + \alpha_A(p_A^L - 2\kappa t_u) - 2t_u(\hat{x}\bar{\alpha} + \alpha_A)((1 - \hat{x})\bar{\alpha} + \alpha_B)$$
 (3)

$$\Pi_B^I = (1 - \hat{x})\bar{\alpha}(p_B - 2\kappa t_u) + \alpha_B(p_B^L - 2\kappa t_u) - 2t_u(\hat{x}\bar{\alpha} + \alpha_A)((1 - \hat{x})\bar{\alpha} + \alpha_B).$$
 (4)

The first term of each function describes a network's direct profits from customers in the battlezone net of Transit costs which stem from sending data to or receiving data from customers of the other network. The second term denotes the same for its locked customers, while the third term adjusts for the traffic that is exchanged between A and B. This term has to be paid to the intermediary by each network, is of equal size for both firms and will become a main formal driver of the model. Note that traffic has to be paid twice for each consumer since we have assumed that all consumers both send data to and receive data from all other consumers. The first-order-condition of network A is given by

$$\frac{\partial \Pi_A^I}{\partial p_A} = \frac{\bar{\alpha}}{2\tau^2} \left( \tau (p_B - 2p_A + \kappa 2t_u - \alpha_A 2t_u + \alpha_B 2t_u + \tau) + p_A 2t_u \bar{\alpha} - p_B 2t_u \bar{\alpha} \right) = 0, \quad (5)$$

while B's is analogous. We derive reaction functions as

$$p_A(p_B) = \frac{2t_u\bar{\alpha} - \tau}{2(t_u\bar{\alpha} - \tau)}p_B + \frac{2\tau t_u(\alpha_A - \alpha_B - \kappa) - \tau^2}{2(t_u\bar{\alpha} - \tau)}$$
(6)

$$p_B(p_A) = \frac{2t_u\bar{\alpha} - \tau}{2(t_u\bar{\alpha} - \tau)}p_A + \frac{2\tau t_u(\alpha_B - \alpha_A - \kappa) - \tau^2}{2(t_u\bar{\alpha} - \tau)}.$$
 (7)

Second-order-conditions are satisfied and the slope of the reaction functions is between zero and one for  $\tau > 2\bar{\alpha}t_u$ , which we assume henceforth. This yields the following equilibrium prices

$$p_A^* = \tau + \frac{2t_u(\kappa(3\tau - 4t_u\bar{\alpha}) - \Delta\tau)}{3\tau - 4t_u\bar{\alpha}} = \tau(1 - z\Delta) + 2\kappa t_u$$
 (8)

$$p_B^* = \tau + \frac{2t_u(\kappa(3\tau - 4t_u\bar{\alpha}) + \Delta\tau)}{3\tau - 4t_u\bar{\alpha}} = \tau(1 + z\Delta) + 2\kappa t_u, \tag{9}$$

where  $\Delta = \alpha_A - \alpha_B \ge 0$  and  $z = \frac{2t_u}{(3\tau - 4t_u\bar{\alpha})} > 0$ . Hence, A's equilibrium market share is

$$\hat{x} = \frac{1}{2} + \frac{2t_u \Delta}{(3\tau - 4t_u \bar{\alpha})} = \frac{1}{2} + z\Delta.$$
 (10)

Equilibrium profits under the Intermediary regime are given by

$$\Pi_A^I = \frac{1}{2}\tau\bar{\alpha}(1 + z\Delta - 2z^2\Delta^2) + z^2\Delta^2\bar{\alpha}(3\tau - 2t_u\bar{\alpha}) - t_uw + \alpha_A(p_A^L - 2\kappa t_u)$$
 (11)

$$\Pi_B^I = \frac{1}{2}\tau\bar{\alpha}(1 - z\Delta - 2z^2\Delta^2) + z^2\Delta^2\bar{\alpha}(3\tau - 2t_u\bar{\alpha}) - t_uw + \alpha_B(p_B^L - 2\kappa t_u).$$
 (12)

It is obvious that A's direct profits from the battlezone,  $\frac{1}{2}\tau\bar{\alpha}(1+z\Delta-2z^2\Delta^2)$ , increase while B's direct profits decrease with growing asymmetry  $\Delta$ . Furthermore, total Transit costs of each network,  $t_uw + \alpha_i 2\kappa t_u - z^2\Delta^2\bar{\alpha}(3\tau - 2t_u\bar{\alpha})$ , are maximized for symmetry ( $\Delta = 0$ ). We find:

**Proposition 1** Under the Intermediary regime of interconnection, network A prices more aggressively than B leading to a higher market share and larger profits of A in the battlezone.

The key to understanding this proposition is that Transit payments of A and B to the intermediary decrease with growing network asymmetry. Thus, the larger network A has larger incentives to increase its market share than the smaller one. If A could sell to a marginal consumer, its income would increase and its Transit costs would decrease. B faces an extra trade-off: if acquiring a marginal customer within the battlezone, its income would increase, but corresponding Transit costs would increase in line. Therefore, A's marginal profit from acquiring another customer is larger than B's making A more aggressive. Similarly, A's ex post profits increase with growing ex ante asymmetry, which also minimizes both networks' Transit payments since more traffic is exchanged "on-net", i.e., if sender and receiver are customers of the same network.

# 3.4 Price Competition under Bill-and-Keep Peering

If networks peer with each other, their profit functions show two differences in relation to the case without Peering: Peering's upside is that networks do not have to pay the intermediary for traffic that is exchanged solely between the two networks involved, anymore. Its downside is that the Peering partners have to set up direct lines, buy new equipment such as routers and have to bear Peering management

costs. All these types of costs are compiled in the variable F, which is not, according to various industry talks, correlated with network size or the amount of traffic transmitted.

This leads to the following profit functions under Peering:

$$\Pi_A^P = \hat{x}\bar{\alpha}(p_A - 2\kappa t_u) + \alpha_A(p_A^L - 2\kappa t_u) - F, \tag{13}$$

$$\Pi_{B}^{P} = (1 - \hat{x})\bar{\alpha}(p_{B} - 2\kappa t_{u}) + \alpha_{B}(p_{B}^{L} - 2\kappa t_{u}) - F. \tag{14}$$

Equilibrium prices can be derived as

$$p_A^* = \tau + 2\kappa t_u \tag{15}$$

$$p_B^* = \tau + 2\kappa t_u, \tag{16}$$

leading to an equilibrium market share for A (and for B, respectively) of

$$\hat{x} = \frac{1}{2}.\tag{17}$$

Equilibrium profits under the Peering regime are given by

$$\Pi_A^P = \frac{1}{2}\tau\bar{\alpha} + \alpha_A(p_A^L - 2\kappa t_u) - F \tag{18}$$

$$\Pi_B^P = \frac{1}{2}\tau\bar{\alpha} + \alpha_B(p_B^L - 2\kappa t_u) - F.$$
 (19)

These equations yield:

**Proposition 2** Under the Peering regime of interconnection, (i) regardless of asymmetries in installed bases networks' pricing behavior is symmetric. (ii) Market shares in the battlezone are symmetric. (iii) Leaving out profits from the installed bases, profits from competition in the battlezone are symmetric. (iv) If installed bases were symmetric ( $\Delta = 0$ ), equilibrium prices and battlezone market shares would be the same under Intermediary and Peering regimes.

The intuition for (i) through (iii) is that, since under a Peering regime Transit costs for traffic between the two parties are waived, the larger network has no extra incentives to undercut the smaller one, anymore. Therefore, incentive structures, behavior and profits are symmetric. This intuition is confirmed by (iv) stating that symmetric networks always behave in the same way regardless of the interconnection regime.

# 4 Bargaining under Paid Peering

Given networks decided to interconnect under the Paid Peering regime, on the second stage of the game we should calculate the settlement-fee, S, one network has to pay the other to make the latter agree to Peering.<sup>20</sup> If they opted for Intermediary or BK, this stage would be waived.

It facilitates further analysis, if we first derive the networks' relative individual incentives to accept Bill-and-Keep Peering.

**Proposition 3** The smaller network always has higher incentives to reach a Bill-and-Keep Peering relative to Intermediary than the larger network.

Proof: Network A's incentives to BK—the gains from Peering—are smaller than B's, if  $\Pi_A^P - \Pi_A^I < \Pi_B^P - \Pi_B^I$ , or (18) – (11) < (19) – (12), which is true for all defined parameter realizations.

Because of Proposition 3, it is clear that network B always has to pay network A under Paid Peering, not vice versa. Let

$$S \equiv \frac{1}{2} (\Pi_B^P - \Pi_B^I - (\Pi_A^P - \Pi_A^I)) = \frac{\bar{\alpha}\tau z\Delta}{2}$$
 (20)

be this settlement B has to pay A, meaning that we assume equal bargaining power and use the respective equilibrium profits under the Intermediary regime as threat points.<sup>21</sup> At a non-cooperative bargaining outcome, the networks share equally any gains relative to their threat points. This formulation ensures that each player obtains (or keeps) profits from the Intermediary case, at least, while only "excess"

 $<sup>^{20}</sup>$ Here, the transfer payment or access charge between networks, unlike in most papers on interconnection in telecommunications, is of a lump-sum type, not a per unit of data fee. The two are structurally similar as long as they do not influence pricing behavior in the retail market. Given that and our assumption of perfect information, S could be interpreted as the sum of all per unit fees in a given period. In contrast, inclusion of access charges that influence retail competition on the third stage of the game is not the focus of our more fundamental paper. Therefore, we follow Besen et al. (2001) in assuming a lump-sum payment.

<sup>&</sup>lt;sup>21</sup>This formulation is analogous to Besen et al. (2001) whose approach is based on the Nash bargaining model of Binmore et al. (1986). It can be applied if we assume that the lack of Peering is sustained only temporarily during bargaining until an agreement is reached, since this resembles the bargaining result according to the non-cooperative bargaining theory with short times between offers.

profits are shared. Therefore, the assumption of equal bargaining power—which is expressed by the factor 1/2 in (20)—is not crucial here since it does not affect absolute incentives to agree to Paid Peering relative to Intermediary.

In general, A's equilibrium profits under Paid Peering are  $\Pi_A^{PP} = \Pi_A^P + S$  while B's are  $\Pi_B^{PP} = \Pi_B^P - S$ . Using (20) yields

$$\Pi_A^{PP} = \frac{1}{2}\tau\bar{\alpha} + \alpha_A(p_A^L - 2\kappa t_u) - F + \frac{\bar{\alpha}\tau z\Delta}{2}$$
(21)

$$\Pi_B^{PP} = \frac{1}{2}\tau\bar{\alpha} + \alpha_B(p_B^L - 2\kappa t_u) - F - \frac{\bar{\alpha}\tau z\Delta}{2}.$$
 (22)

# 5 Regime Equilibria

Being aware of Nash equilibria in prices given the respective regimes, we now proceed to analyze incentives on the first stage: When do networks wish to peer with a specific competitor? What form of Peering would prevail if side payments were feasible?

Before analyzing equilibria, we are to specify the support of  $\Delta$  in general. (10) explicates that to receive interior solutions for  $\hat{x}$  so that  $\hat{x} \in [0, 1]$ , it is necessary that  $\Delta \in [-\frac{1}{2z}, \frac{1}{2z}]$ . Thus, as  $\Delta \geq 0$  by definition, we have  $\Delta_{max} \equiv \frac{1}{2z}$ , where  $t_u \geq \frac{3\tau}{4(1+\bar{\alpha})}$  which is always true for defined values. If  $\Delta$  lies outside of these boundaries, the larger network's aggressiveness in the price competition is so strong that the smaller network will be driven out of the (battlezone) market. Henceforth, we restrict our analysis to  $\Delta \in [0, \Delta_{max}]$ .

Now, recall that each network can force the other one to play the Intermediary strategy. If and only if both parties either agree on BK or on PP, that regime will be an equilibrium. Therefore, following our assumption that the Intermediary regime is the status-quo when playing the first stage of the game, Intermediary is a Nash equilibrium for all levels of asymmetry. This might explain why we observe usage of IP-Transit among both symmetric and asymmetric networks in practice, given the cost of alternative Peering regimes, F, are not too low.

When is Bill-and-Keep Peering a candidate for equilibrium outcomes? It is an equilibrium strategy for both A and B if no player has an incentive to deviate from it and to obtain the Intermediary outcome. We find

**Proposition 4** Bill-and-Keep Peering is an equilibrium outcome for small network asymmetries (where  $\Delta \in [0, \Delta_{BK}]$ ).

Proof: refer to the appendix.

When is Paid Peering a candidate for equilibrium outcomes? Because of our assumption that under Paid Peering "excessive" joint profits can be perfectly exchanged, no player will prefer the Intermediary strategy over PP as long as  $\sum \Pi_i^{PP} > \sum \Pi_i^I$ . Rearranging this relation based on equations (11), (12), (21) and (22) yields that Paid Peering is an equilibrium outcome if

$$\Delta < \sqrt{\frac{t_u w - F}{2z^2 \bar{\alpha}(\tau - t_u \bar{\alpha})}} \equiv \Delta_P. \tag{23}$$

According to our assumptions, we always have  $\Delta_P \geq 0$ . Via resubstitution of z we find that  $\Delta_P < \Delta_{max}$  for  $t_u w - F < \frac{\bar{\alpha}}{2}(\tau - t_u \bar{\alpha})$ . Summarizing, if F is sufficiently low  $(F \leq t_u w)$ , networks' interconnection decision is largely dependent on their ex ante size asymmetry: Paid Peering dominates Intermediary for low  $\Delta$ , and vice versa for large  $\Delta$ . But if F is too low, Intermediary can never be an equilibrium as (23) emphasizes that networks will interconnect via an Intermediary if the difference in size of two networks is relatively large.

**Proposition 5** Assume  $t_u w - \frac{\bar{\alpha}}{2}(\tau - t_u \bar{\alpha}) < F < t_u w$ . (i) Intermediary constitutes a Nash equilibrium for any level of network asymmetry. (ii) However, for low asymmetry networks prefer Peerings: (a) For all  $\Delta \in [0, \Delta_{BK}]$ , both Bill-and-Keep and Paid Peering form equilibrium outcomes. (b) For all  $\Delta \in (\Delta_{BK}, \Delta_P)$ , Paid Peering is an equilibrium outcome. (iii) For all  $\Delta \in (\Delta_P, \Delta_{max}]$ , Intermediary is a unique equilibrium outcome. (iv) For  $\Delta = \Delta_P$ , both Paid Peering and Intermediary are equilibrium outcomes.

Proof: refer to the appendix.

Figure 2 provides a graphical intuition of Proposition 5: It plots where (a) Bill-and-Keep Peering, (b) Paid Peering and (c) Intermediary constitute Nash equilibria. Left of  $\Delta_P$ , at least one Peering regime is preferred by the networks over buying IP-Transit—and they can deviate from playing an Intermediary strategy without risk. We have used a dashed line to indicate this.

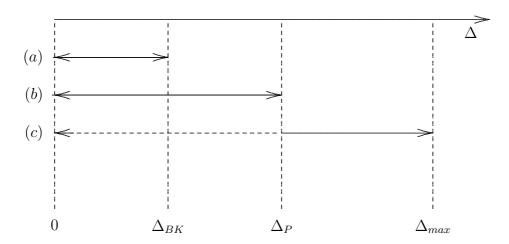


Figure 2: First stage Nash equilibria

Amongst others, this proposition suggests an intuition why, according to anecdotal evidence, Bill-and-Keep Peering has been the dominant Peering regime in practice. If networks are sufficiently symmetric and the smaller network can credibly announce that it will not bargain over a settlement-fee, the larger network is better off by accepting BK instead of being tough, too, and ending up paying the Intermediary.<sup>22</sup> If networks' asymmetry is small but not very small, the smaller network knows that the larger one would never accept BK because the Intermediary outside option is more attractive. Then, the smaller network is better off by paying some of its gains from Peering via a settlement-fee thereby compensating the larger one for its losses. Reflecting on these two arguments indicates that in practice—and outside of our model—the sequence of moves is crucial.

<sup>&</sup>lt;sup>22</sup>One reason for the smaller network's resistance to bargain at all could be explained by the fact that the bargaining process associated with Paid Peering may involve extra transaction costs in comparison to BK. Another explanation could be *legacy* which is, however, questionable from an economic point of view. The argument claims that, at the beginning of the commercial internet era, networks did not focus on the strategic aspects of interconnection but strived for reaching world-wide connectivity fast. Nowadays, they found themselves in the resource consuming process of reviewing their existing Peering policies.

### 6 Welfare

Now we know which interconnection regime networks will choose given exogenous parameter realizations. But are market outcomes beneficial for consumers and total welfare, as well?

### 6.1 Consumer Surplus

We restrict the analysis to the  $\bar{\alpha}$  consumers residing in the battlezone since consumer surplus within the locked regions is neither a function of the networks' interconnection regime nor of their battlezone prices. Hence aggregate consumer surplus is the integral over individual net benefit (according to (1)) using the marginal consumer as boundary. As under (Paid) Peering, equilibrium prices of networks A and B are equal and each one gets a market share of 0.5, we can calculate consumer surplus as

$$CS^{P} = 2\bar{\alpha} \int_{0}^{0.5} (v - \tau x - p_{A}) dx = \bar{\alpha} (v - \frac{5}{4}\tau - 2\kappa t_{u}).$$
 (24)

In contrast, consumer surplus under Intermediary is denoted by

$$CS^{I} = \bar{\alpha} \left( \int_{0}^{\hat{x}} (v - \tau x - p_{A}) dx + \int_{\hat{x}}^{1} (v - \tau (1 - x) - p_{B}) dx \right)$$
$$= \bar{\alpha} \left( v - \frac{5}{4} \tau - 2\kappa t_{u} \right) + \bar{\alpha} \tau z^{2} \Delta^{2} = CS^{P} + \bar{\alpha} \tau z^{2} \Delta^{2}. \quad (25)$$

Analogous to section 3,  $CS^P = CS^I$  if networks are symmetric ( $\Delta = 0$ ). But for all  $\Delta > 0$  consumer surplus is larger under the Intermediary regime. This is intuitive since in the Intermediary case the larger network competes more aggressively in prices than in the Peering case, but it also obtains a higher market share within the battlezone. Hence a majority of consumers enjoys extra surplus which is not offset completely by higher prices that are paid by the fewer customers of the smaller network. It is straightforward to observe from (25) that consumer surplus under Intermediary relative to Peering increases even further with growing network asymmetry.

### 6.2 Total Welfare

Up to which asymmetry should networks peer from a social perspective? Clearly, we can find this point,  $\Delta_P^{Soc}$ , where a social planner including both consumer surplus and

producer surplus (i.e. profits of networks A and B and the intermediary network) in his calculation would be indifferent between Peering and Intermediary. As  $\sum \Pi_i^P = \sum \Pi_i^{PP}$ , we can find this level via setting

$$CS^{P} + \Pi_{A}^{P} + \Pi_{B}^{P} + \Pi_{Int}^{P} = CS^{I} + \Pi_{A}^{I} + \Pi_{B}^{I} + \Pi_{Int}^{I}$$
(26)

where profits of the intermediary are denoted by  $\Pi_{Int}^P = 2\kappa t_u(\alpha_A + \alpha_B + \bar{\alpha})$  and  $\Pi_{Int}^I = \Pi_{Int}^P + 2t_u w - 2z^2 \Delta^2 \bar{\alpha} (3\tau - 2t_u \bar{\alpha})$  respectively. Employing equations (24), (18) and (19) as well as (25), (11) and (12) yields that from a social perspective networks should peer if

$$\Delta > \sqrt{\frac{2F}{z^2 \tau \bar{\alpha}}} \equiv \Delta_P^{Soc}. \tag{27}$$

However, since currently all major intermediary backbones are US based firms,<sup>23</sup> one might also be interested in the ranges of asymmetry where a non-US policy maker would like networks to peer, i.e. without taking into account the profits of the intermediary network. Therefore, we set

$$CS^{P} + \Pi_{A}^{P} + \Pi_{B}^{P} = CS^{I} + \Pi_{A}^{I} + \Pi_{B}^{I}$$
(28)

and find that in this "trade policy" case, a regulator would want networks to peer as long as

$$\Delta < \sqrt{\frac{2(t_u w - F)}{z^2 \bar{\alpha}(5\tau - 4t_u \bar{\alpha})}} \equiv \Delta_P^{TP}.$$
 (29)

It might be startling that both a trade policy regulator and the profit maximizing networks prefer Peering for a lesser degree of asymmetry, while a social planner prefers Peering for larger asymmetry. To understand the intuition of the three  $\Delta$ -thresholds recall that the respective optimizers include different parameters in their calculi.

Networks trade-off Peering costs F versus Transit costs depending on  $t_u$ . If  $\Delta$  increases, F remains constant while joint Transit costs decrease. Therefore, above a certain level of asymmetry,  $\Delta_P$ , networks prefer the Intermediary regime.

A "trade policy" regulator faces the same trade-off and hence prefers Peering for low levels of asymmetry. But in addition he regards consumer surplus, which grows with  $\Delta$  under Intermediary due to fiercer network competition but remains constant

<sup>&</sup>lt;sup>23</sup>See http://www.fixedorbit.com/stats.htm.

under Peering. Therefore, trade policy makers wish to have the Intermediary regime implemented even for lower levels of asymmetry than networks themselves.

A social planner, in contrast, does not observe the effect of decreasing Transit costs for larger asymmetry as this money flows to the intermediary backbone, which is included in his optimization calculus. Therefore, for low levels for asymmetry he only takes into account Peering costs F and prefers Intermediary regimes. With rising  $\Delta$ , under Intermediary the social planner observes distortions due to networks' fiercer competition, which depend on the transportation cost  $\tau$  in the model. As a consequence, above a threshold,  $\Delta_P^{Soc}$ , he prefers interconnection via Peering regimes.

**Proposition 6** (i) Excess Peering: The level of asymmetry of network sizes up to which a "trade policy" regulator would prefer Peering,  $\Delta_P^{TP}$ , is smaller than the asymmetry up to which networks peer without regarding consumer welfare,  $\Delta_P$ . (ii) Within the range where networks peer but where it is suboptimal from a "trade policy" viewpoint, the loss increases with growing asymmetry.

Proof: see appendix.

Now we know that always  $\Delta_P^{TP} < \Delta_P$ . However,  $\Delta_P^{Soc}$  is not fixed within this range. What happens for low, medium and large realizations of  $\Delta_P^{Soc}$ , and when do those cases occur? We distinguish among three possible realizations. Please, recall that the minimum level of F is  $wt_u - \frac{\bar{\alpha}}{2}(\tau - t_u\bar{\alpha})$  and its maximum level is  $wt_u$ :

- Case I:  $\Delta_P^{Soc} \leq \Delta_P^{TP} \quad \forall \quad F \in (wt_u \frac{\bar{\alpha}}{2}(\tau t_u\bar{\alpha}), \frac{\tau}{6\tau 4t_u\bar{\alpha}}wt_u]$
- Case II:  $\Delta_P^{TP} < \Delta_P^{Soc} \le \Delta_P \quad \forall \quad F \in (\frac{\tau}{6\tau 4t_u\bar{\alpha}}wt_u, \frac{\tau}{5\tau 4t_u\bar{\alpha}}wt_u]$
- Case III:  $\Delta_P < \Delta_P^{Soc} \quad \forall \quad F \in \left(\frac{\tau}{5\tau 4t_u\bar{\alpha}}wt_u, wt_u\right)$

By checking these cases with the respective definitions of  $\Delta_P$ ,  $\Delta_P^{Soc}$  and  $\Delta_P^{TP}$ , we easily observe

Proposition 7 (i) Within cases I and III but not in case II, there exist ranges where the equilibrium interconnection regime is in line with the views of both a social planner and a "trade policy" regulator. In case I (III) Peering (Intermediary) is optimal from these three perspectives as long as  $\Delta_P^{Soc} < \Delta < \Delta_P^{TP}$  ( $\Delta_P < \Delta < \Delta_P^{SOC}$ ). (ii) If Peering costs are sufficiently large, Peering never occurs where it is socially efficient ( $\Delta_P < \Delta_P^{SOC}$ ). (iii) If Peering costs are sufficiently large, "trade policy regulators" only support Peering where it is socially inefficient ( $\Delta_P^{TP} < \Delta_P^{SOC}$ ).

Proof: see appendix.

Figure 3 provides a graphical intuition for Propositions 6 and 7. Peering is preferred by (a) the networks themselves (b) a "trade policy" regulator (c) a total welfare maximizer.

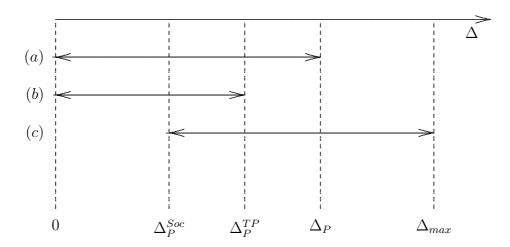


Figure 3: Welfare considerations for  $\Delta_P^{Soc} \leq \Delta_P^{TP}$  (Case I)

Therefore, it is possible that both types of regulators are content with networks' actions, but it is also feasible that they would like to intervene into the market. We can characterize as a general rule that networks always peer excessively from a "trade policy" regulator's point of view.

### 7 Discussion

We verified three possible sources of revenues for IBPs: end-user charges  $(p_i)$ , Paid Peering revenues (S) and IP-Transit fees  $(t_u)$ .

Given that IP-Transit is a homogenous good, as long as this market is characterized by (i) perfect competition, (ii) the absence of bottlenecks and (iii) excess capacity, we should expect to observe Transit charges declining to marginal costs  $(t_u = MC = 0)$ . This view is supported by empirical data: OECD (2002) states that prices for IP-Transit had fallen by up to 55 percent annually from 1998-2000. Furthermore, Telegeography notes that "in 2004, backbone access prices around the

world fell about 50 percent over the previous year. This year [2005] prices fell between 23 and 33 percent."<sup>24</sup>

As a consequence of this development, we should expect revenues from IP-Transit to vanish. Moreover, according to (20) settlement-fees in case of Paid Peering should follow the direction of Transit fees, i.e., they should approach zero, too. Thus, profits from Bill-and-Keep and Paid Peering would converge. (23) shows that as a boundary solution we obtain  $\Delta_P \to 0$  meaning that Intermediary would become the dominant interconnection regime.<sup>25</sup> Similarly, we find that  $\Delta_{max} \to (+\infty)$  encouraging market entry by small IBPs.

This leaves networks with equilibrium profits of  $\Pi_i = \frac{1}{2}\tau\bar{\alpha} + \alpha_i p_i^L$ , according to (11) and (12). These profits heavily depend on  $\tau$ , i.e., IBPs should exert any effort to differentiate themselves in the retail market, for instance by introducing new, specific services for end-users.<sup>26</sup>

# 8 Conclusion

Our main objective was to study IBPs' optimal interconnection decisions, which are strategically linked to competition for end-users. Based on our propositions we find the following main practical implications:

- 1. If, besides Intermediary and Bill-and-Keep, networks also consider Paid Peering as a possible type of interconnection, we expect to observe more Paid Peering in the future. This translates to more Peering agreements in general which, in turn, leads to higher profits of IBPs.
- 2. This development harms consumer surplus, however.
- 3. Since the emergence of Paid Peering also lowers demand for IP-Transit, top level backbones will lose revenues.
- 4. As all top level backbones are US-based, non-US policy makers do not include profits from IP-Transit in their calculus. Instead of considering to punish

<sup>&</sup>lt;sup>24</sup>See http://www.telegeography.com/press/releases/2005-08-23.php.

<sup>&</sup>lt;sup>25</sup>Technically,  $\Delta_P \to (-\infty)$ .

 $<sup>^{26}</sup>$ See Prüfer and Jahn (2005) for more empirical data and a discussion of the internet backbone market's structure and development.

large networks who refuse (Bill-and-Keep) Peering to smaller ones, these policy makers should review to restrict Peering, since networks do not care about the fact that fiercer competition under Intermediary benefits consumers, and peer excessively. In contrast, since US-based policy makers do account for profits from IP-Transit, they favor Peerings among networks sufficiently asymmetric in size. Hence, they should seek to discourage large networks from refusing to peer with smaller ones.

These results could also be applied to a telecommunications market which was both unregulated in terms of inter-carrier compensation fees and not subject to price discrimination regarding destinations of calls.

Since this paper adopts a new approach to network interconnection on the internet, it could be extended in various directions in future research. One option is to analyze thoroughly the implications of different Paid Peering contracts, e.g. an ex ante lump-sum settlement (as assumed here) vs. ex post payment of a price  $p^P < t_u$  per unit of data that influences networks' competitive retail pricing. Another interesting issue would be to study the impact of one (or both) networks being a tier-1 network on our analysis. This could include endogenization of the Transit charge.

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# A Appendix

# A.1 Proof of Proposition 4

If the larger network A prefers BK over Intermediary, according to Proposition 3 the smaller network B will do so as well. Hence we can prove Proposition 4 by proving existence of a defined parameter range where  $\pi_A^P > \pi_A^I$ . As we assumed  $F \leq t_u w$ ,  $\pi_A^P > \pi_A^I$  holds for  $\Delta = 0$ , which forms the lower boundary of this range. Define  $\Delta_{BK}$  as the upper boundary. As long as  $t_u w - F \geq \frac{1}{4}\bar{\alpha}(3\tau - 2t_u\bar{\alpha})$  we have  $\Delta_{BK} = \Delta_{max}$  as a corner solution. As a consequence,  $\pi_A^P > \pi_A^I$  holds for all defined values of  $\Delta$ . For  $t_u w - F < \frac{1}{4}\bar{\alpha}(3\tau - 2t_u\bar{\alpha})$ ,  $\Delta_{BK} < \Delta_{max}$ . Then,  $\pi_A^P < \pi_A^I$  at  $\Delta_{max}$ . According to equation (18),  $\pi_A^P(\Delta)$  is constant. According to (11),  $\pi_A^I(\Delta)$  is a continuous, strictly increasing function on  $[0, \Delta_{max}]$ . Therefore,  $\Delta_{BK}$  exists and is unique. Thus, for

 $\Delta \in [0, \Delta_{BK}]$  both networks will not deviate from a BK strategy, given the other party does not deviate.  $\square$ 

### A.2 Proof of Proposition 5

Proof: Ad (i): This follows from our assumption that the agreement of both networks is needed to deviate from the Intermediary regime.

Ad (ii.a): For  $\Delta = \Delta_{BK}$ , by definition we have

$$\pi_A^P = \pi_A^I, \tag{A.1}$$

and, by Proposition 3, there we have

$$\pi_B^P > \pi_B^I. \tag{A.2}$$

We shall distinguish among three cases:

- 1. Assume  $\Delta_{BK} = \Delta_P$ . Then,  $\Delta_P$  requires  $\pi_A^P + \pi_B^P = \pi_A^I + \pi_B^I$ . Substituting (A.1) in this condition yields  $\pi_B^P = \pi_B^I$ , which is in contradiction to (A.2).
- 2. Assume  $\Delta_{BK} > \Delta_P$ . Then,  $\Delta_P$  requires  $\pi_A^P + \pi_B^P < \pi_A^I + \pi_B^I$ . Substituting (A.1) in this condition yields  $\pi_B^P < \pi_B^I$ , which is in contradiction to (A.2).
- 3. Assume  $\Delta_{BK} < \Delta_P$ . Then,  $\Delta_P$  requires  $\pi_A^P + \pi_B^P > \pi_A^I + \pi_B^I$ . Substituting (A.1) in this condition yields  $\pi_B^P > \pi_B^I$ , which is in line with (A.2).

Therefore, for  $\Delta \in [0, \Delta_{BK}]$  both BK and PP are equilibria, while for  $\Delta \in (\Delta_{BK}, \Delta_P)$  PP is a unique equilibrium. (ii.b), (iii) and (iv) follow from our above argumentation.  $\square$ 

# A.3 Proof of Proposition 6

Ad~(i):  $\Delta_P$  and  $\Delta_P^{TP}$  both have the same denominator. Therefore  $\Delta_P^{TP} < \Delta_P$  if  $2\bar{\alpha}t_u^2(5\tau - 4t_u\bar{\alpha}) > 8\bar{\alpha}t_u^2(\tau - t_u\bar{\alpha})$ , which is true for all defined parameter realisations. Ad~(ii): The loss (L) accumulates to

$$L = \{CS^I(\Delta^2) + \sum \prod_i^I(\Delta^2) - (CS^P + \sum \prod_i^P) | \Delta_P^{TP} \le \Delta \le \Delta_P \}.$$
 (A.3)

By using (25) and the fact that  $\frac{\partial}{\partial \Delta^2} (\bar{\alpha} \tau z^2 \Delta^2) > 0$  and  $\frac{\partial}{\partial \Delta^2} (\sum \Pi_i^I(\Delta^2)) > 0$ , it follows that

$$\frac{\partial L}{\partial \Delta^2} > 0. \quad \Box$$
 (A.4)

# A.4 Proof of Proposition 7

Ad (i): This follows directly from the respective definitions.

Ad (ii): Peering occurs if  $\Delta < \Delta_P$ . It is efficient if  $\Delta > \Delta_P^{Soc}$ . It never occurs when it is efficient if  $\Delta_P < \Delta_P^{Soc}$ . This is true for all  $F \in (\frac{\tau}{5\tau - 4t_u\alpha_c}wt_u, wt_u)$ .

Ad (iii): A "trade policy" regulator supports Peering if  $\Delta < \Delta_P^{TP}$ . Peering is efficient if  $\Delta > \Delta_P^{Soc}$ . It is never supported by a "trade policy" regulator when it is efficient if  $\Delta_P^{TP} < \Delta_P^{Soc}$ . This is true for all  $F \in (\frac{\tau}{6\tau - 4t_u\alpha_c}wt_u, wt_u)$ .